

Graph Connectivity in Sparse Subspace Clustering

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- 1 Subspace Clustering
- 2 Sparse Subspace Clustering
- 3 Graph Connectivity
- 4 Conclusion

- 1 Subspace Clustering
 - Example
 - Subspace Clustering
 - Applications
 - Solutions
- 2 Sparse Subspace Clustering
- 3 Graph Connectivity
- 4 Conclusion

Frame 1:



- $\mathbf{x}_i = [A_1]_{2 \times 4} \begin{bmatrix} \mathbf{x}_i \\ 1 \end{bmatrix}_{4 \times 1}$

Frame 1:



- $\mathbf{x}_j = [A_1]_{2 \times 4} \begin{bmatrix} \mathbf{x}_j \\ 1 \end{bmatrix}_{4 \times 1}$
- $[\mathbf{x}_1 \cdots \mathbf{x}_n] = [A_1]_{2 \times 4} \begin{bmatrix} \mathbf{x}_1 \cdots \mathbf{x}_n \\ 1 \cdots 1 \end{bmatrix}_{4 \times n}$

Frame 1:



Frame 2:



- $\mathbf{x}_i = [A_1]_{2 \times 4} \begin{bmatrix} \mathbf{x}_i \\ 1 \end{bmatrix}_{4 \times 1}$
- $[\mathbf{x}_1 \cdots \mathbf{x}_n] = [A_1]_{2 \times 4} \begin{bmatrix} \mathbf{x}_1 \cdots \mathbf{x}_n \\ 1 \cdots 1 \end{bmatrix}_{4 \times n}$

Frame 1:



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Frame 2:



- $\mathbf{y}_i = [A_2]_{2 \times 4} \begin{bmatrix} \mathbf{x}_i \\ 1 \end{bmatrix}_{4 \times 1}$

Frame 1:



- $\mathbf{x}_i = [A_1]_{2 \times 4} \begin{bmatrix} \mathbf{X}_i \\ 1 \end{bmatrix}_{4 \times 1}$
- $[\mathbf{x}_1 \cdots \mathbf{x}_n] = [A_1]_{2 \times 4} \begin{bmatrix} \mathbf{x}_1 \cdots \mathbf{x}_n \\ 1 \cdots 1 \end{bmatrix}_{4 \times n}$

Frame 2:



- $\mathbf{y}_i = [A_2]_{2 \times 4} \begin{bmatrix} \mathbf{X}_i \\ 1 \end{bmatrix}_{4 \times 1}$
- $[\mathbf{y}_1 \cdots \mathbf{y}_n] = [A_2]_{2 \times 4} \begin{bmatrix} \mathbf{x}_1 \cdots \mathbf{x}_n \\ 1 \cdots 1 \end{bmatrix}_{4 \times n}$

Frame 1:



- $\mathbf{x}_j = [A_1]_{2 \times 4} \begin{bmatrix} \mathbf{X}_j \\ 1 \end{bmatrix}_{4 \times 1}$
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Frame 2:



- $\mathbf{y}_i = [A_2]_{2 \times 4} \begin{bmatrix} \mathbf{X}_i \\ 1 \end{bmatrix}_{4 \times 1}$
- $\begin{bmatrix} \mathbf{x}_1 \cdots \mathbf{x}_n \\ \mathbf{y}_1 \cdots \mathbf{y}_n \end{bmatrix} = \begin{bmatrix} A_1 \\ A_2 \end{bmatrix}_{4 \times 4} \begin{bmatrix} \mathbf{X}_1 \cdots \mathbf{X}_n \\ 1 \cdots 1 \end{bmatrix}_{4 \times n}$

Frame 1:



- $\mathbf{x}_j = [A_1]_{2 \times 4} \begin{bmatrix} \mathbf{X}_j \\ 1 \end{bmatrix}_{4 \times 1}$
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Frame 2:



- $\mathbf{y}_i = [A_2]_{2 \times 4} \begin{bmatrix} \mathbf{X}_i \\ 1 \end{bmatrix}_{4 \times 1}$
- $\begin{bmatrix} \mathbf{x}_1 \cdots \mathbf{x}_n \\ \mathbf{y}_1 \cdots \mathbf{y}_n \\ \mathbf{z}_1 \cdots \mathbf{z}_n \end{bmatrix} = \begin{bmatrix} A_1 \\ A_2 \\ A_3 \end{bmatrix}_{6 \times 4} \begin{bmatrix} \mathbf{X}_1 \cdots \mathbf{X}_n \\ 1 \cdots 1 \end{bmatrix}_{4 \times n}$

Frame 1:



- $\mathbf{x}_i = [A_1]_{2 \times 4} \begin{bmatrix} \mathbf{x}_i \\ 1 \end{bmatrix}_{4 \times 1}$
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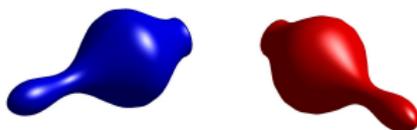
Frame 2:



- $\mathbf{y}_i = [A_2]_{2 \times 4} \begin{bmatrix} \mathbf{x}_i \\ 1 \end{bmatrix}_{4 \times 1}$
- $\begin{bmatrix} \mathbf{x}_1 \cdots \mathbf{x}_n \\ \mathbf{y}_1 \cdots \mathbf{y}_n \\ \vdots \\ \mathbf{z}_1 \cdots \mathbf{z}_n \end{bmatrix} = \begin{bmatrix} A_1 \\ A_2 \\ \vdots \\ A_F \end{bmatrix}_{2F \times 4} \begin{bmatrix} \mathbf{x}_1 \cdots \mathbf{x}_n \\ 1 \cdots 1 \end{bmatrix}_{4 \times n}$

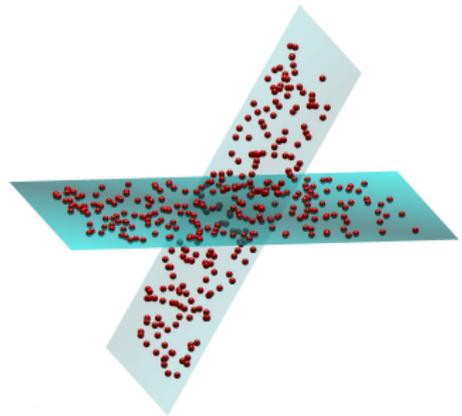
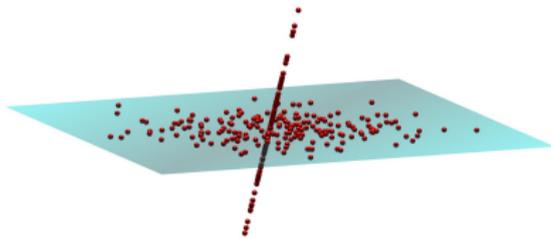


$$\begin{bmatrix} \mathbf{x}_{11} \cdots \mathbf{x}_{1n} \\ \mathbf{x}_{21} \cdots \mathbf{x}_{2n} \\ \vdots \\ \mathbf{x}_{F1} \cdots \mathbf{x}_{Fn} \end{bmatrix}_{2F \times n} = \begin{bmatrix} A_1 \\ A_2 \\ \vdots \\ A_F \end{bmatrix}_{2F \times 4} \begin{bmatrix} \mathbf{x}_1 \cdots \mathbf{x}_n \\ 1 \cdots 1 \end{bmatrix}_{4 \times n}$$

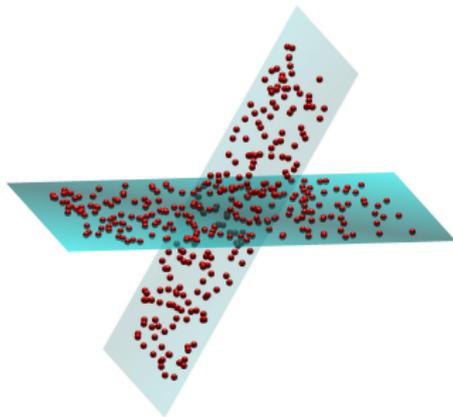
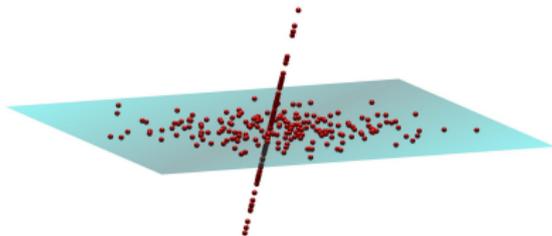


$$\begin{bmatrix} \mathbf{x}_{11} \cdots \mathbf{x}_{1n} \\ \mathbf{x}_{21} \cdots \mathbf{x}_{2n} \\ \vdots \\ \mathbf{x}_{F1} \cdots \mathbf{x}_{Fn} \end{bmatrix}_{2F \times n} = ?$$

- Data lying on a mixture of subspaces.

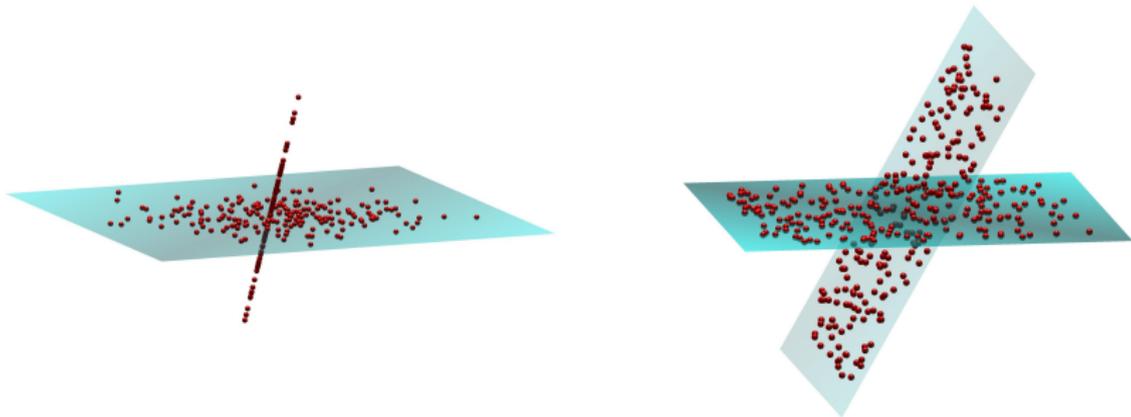


- Data lying on a mixture of subspaces.



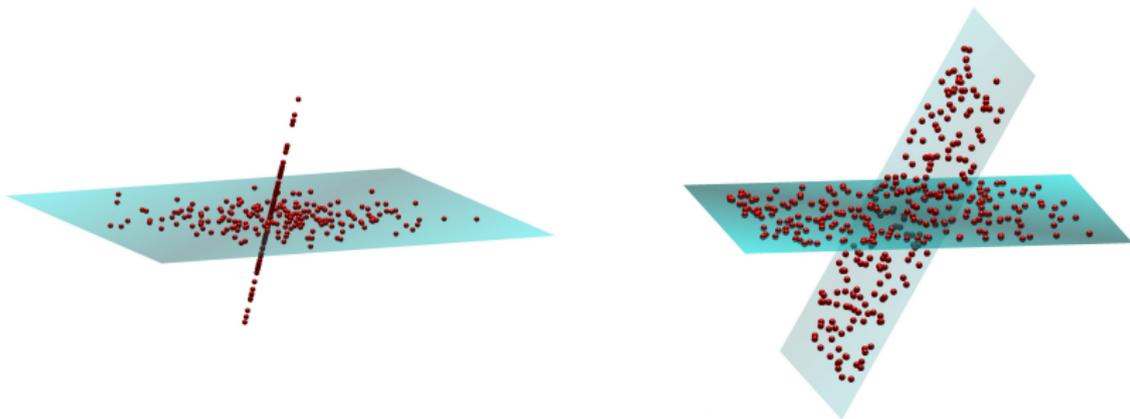
- Number of subspaces

- Data lying on a mixture of subspaces.



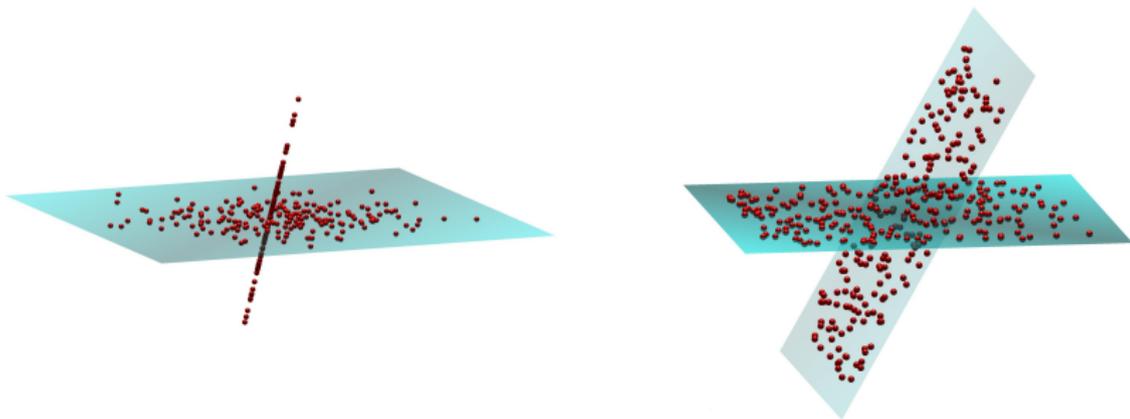
- No. of subspaces + their dimensions

- Data lying on a mixture of subspaces.



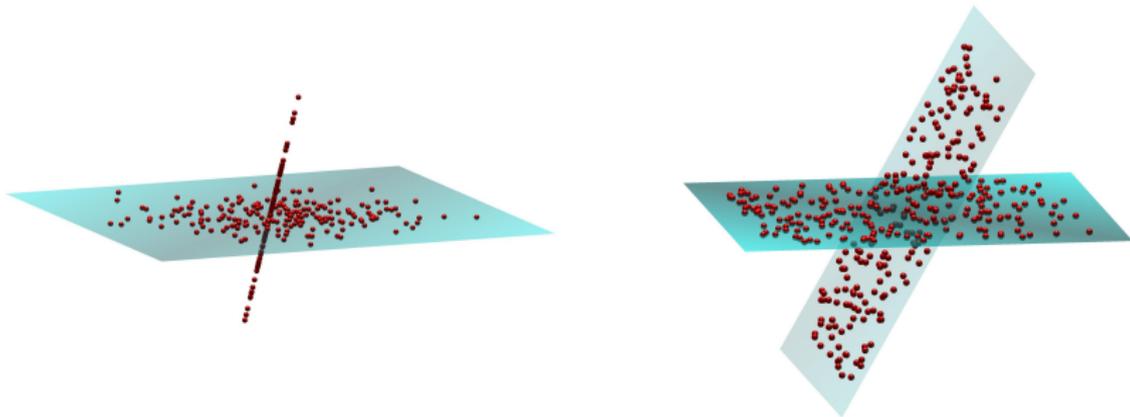
- No. of subspaces + dimensions + A basis for each subspace

- Data lying on a mixture of subspaces.



- No. of subspaces + dimensions + bases + data segmentation

- Data lying on a mixture of subspaces.



- No. of subspaces + dimensions + bases + segmentation

- Motion Segmentation [*Rene Vidal et al. 2008*]
- Video Shot Segmentation [*Le Lu and R. Vidal 2006*]
- Illumination Invariant Clustering [*J. Ho et al. 2003*]
- Image Segmentation [*Alen Yang et al. 2008*]
- Image Representation and Compression [*Wei Hong et al. 2005*]
- Linear Hybrid Systems Identification [*Rene Vidal et al. 2003*]

- Random Sample Consensus (RANSAC) [Martin Fischler and R. Bolles 1981]
- Mixture of Probabilistic PCA [Michael Tipping and C. Bishop 1999]
- Generalized PCA (GPCA) [Rene Vidal et al. 2005]
- Locally Linear Manifold Clustering (LLMC) [Alvina Goh and R. Vidal 2007]
- Agglomerative Lossy Compression (ALC) [Yi Ma et al. 2007]
- Sparse Subspace Clustering (SSC) [Ehsan Elhamifar and R. Vidal 2009]
- Low-Rank Subspace Clustering (LLR) [Guangcan Li et al. 2010]

- 1 Subspace Clustering
- 2 **Sparse Subspace Clustering**
 - Sparse Representation
 - Main Theorem
 - Noise and Outliers
 - Open Problems
- 3 Graph Connectivity
- 4 Conclusion

- Represent \mathbf{y} as a linear combination of the **smallest possible subset** of the vectors $\{\mathbf{x}_1, \mathbf{x}_2, \dots, \mathbf{x}_n\}$.

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$$\min \|\mathbf{a}\|_0 \quad \text{s.t.} \quad \mathbf{y} = \mathbf{X}\mathbf{a}$$

where $\mathbf{X} = [\mathbf{x}_1 \mathbf{x}_2 \cdots \mathbf{x}_n]$.

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- **NP-hard**

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- **NP-hard**
- Use L^1 -minimization:

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- Represent \mathbf{y} as a linear combination of the **smallest possible subset** of the vectors $\{\mathbf{x}_1, \mathbf{x}_2, \dots, \mathbf{x}_n\}$.

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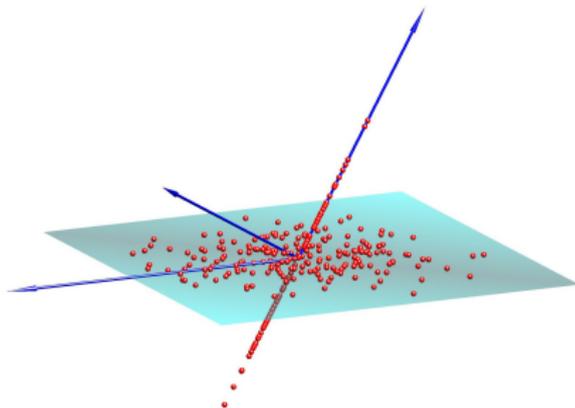
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- **NP-hard**
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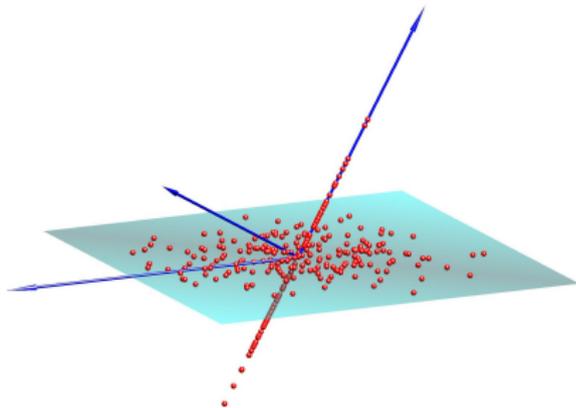
$$\min \|\mathbf{a}\|_1 \quad \text{s.t.} \quad \mathbf{y} = \mathbf{X}\mathbf{a}$$

- L^1/L^0 equivalence

- If we had the basis $[\mathbf{b}_1 \mathbf{b}_2 \dots \mathbf{b}_m]$

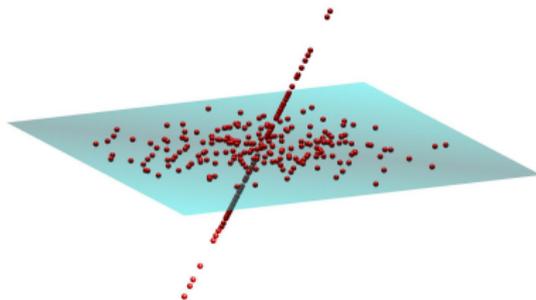


- If we had the basis $[\mathbf{b}_1 \mathbf{b}_2 \dots \mathbf{b}_m]$

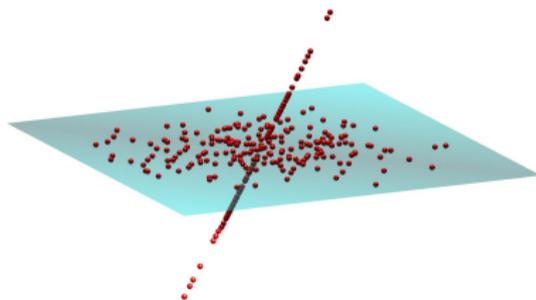


- Represent each \mathbf{x}_i as a sparse combination of $\{\mathbf{b}_1 \mathbf{b}_2 \dots \mathbf{b}_m\}$

- We do not have the basis



- We do not have the basis



- Represent each \mathbf{x}_i as a sparse combination of $X - \{\mathbf{x}_i\}$

- For each \mathbf{x}_j solve:

$$\mathbf{a}^i = \arg \min \|\mathbf{a}\|_1 \quad \text{s.t.} \quad \mathbf{x}_j = \mathbf{X} \mathbf{a}, \quad a_j = 0$$

- For each \mathbf{x}_i solve:

$$\mathbf{a}^i = \arg \min \|\mathbf{a}\|_1 \quad \text{s.t.} \quad \mathbf{x}_i = \mathbf{X} \mathbf{a}, \quad a_j = 0$$

- Construct a graph whose nodes are \mathbf{x}_i and each node \mathbf{x}_i is connected to the node \mathbf{x}_j if the j -th element of \mathbf{a}^i is nonzero.

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Theorem

If the subspaces are independent, the graphs of different subspaces are disconnected.

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Theorem

If the subspaces are independent, the graphs of different subspaces are disconnected.

- Find the **connected components**

- For each \mathbf{x}_i solve:

$$\mathbf{a}^i = \arg \min \|\mathbf{a}\|_1 \quad \text{s.t.} \quad \mathbf{x}_i = \mathbf{X} \mathbf{a}, \quad a_j = 0$$

- Construct a graph whose nodes are \mathbf{x}_i and each node \mathbf{x}_i is connected to the node \mathbf{x}_j if the j -th element of \mathbf{a}^i is nonzero.

Theorem

If the subspaces are independent, the graphs of different subspaces are disconnected.

- Find the **connected components**
- In practice **spectral clustering** using $[\mathbf{a}^1 \mathbf{a}^2 \dots \mathbf{a}^n]$

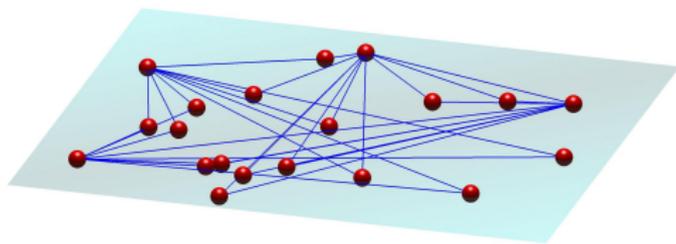


Figure: An example of an SSC graph

- Noise

$$\min \|\mathbf{a}\|_1 + \lambda \|\mathbf{e}\|_2$$

s.t.

$$\mathbf{x}_j = \mathbf{X}_{-j}\mathbf{a} + \mathbf{e}$$

$$a_j = 0$$

- Noise

$$\begin{aligned} \min & \| \mathbf{a} \|_1 + \lambda \| \mathbf{e} \|_2 \\ \text{s.t.} & \\ & \mathbf{x}_i = \mathbf{X}_{-i} \mathbf{a} + \mathbf{e} \\ & a_j = 0 \end{aligned}$$

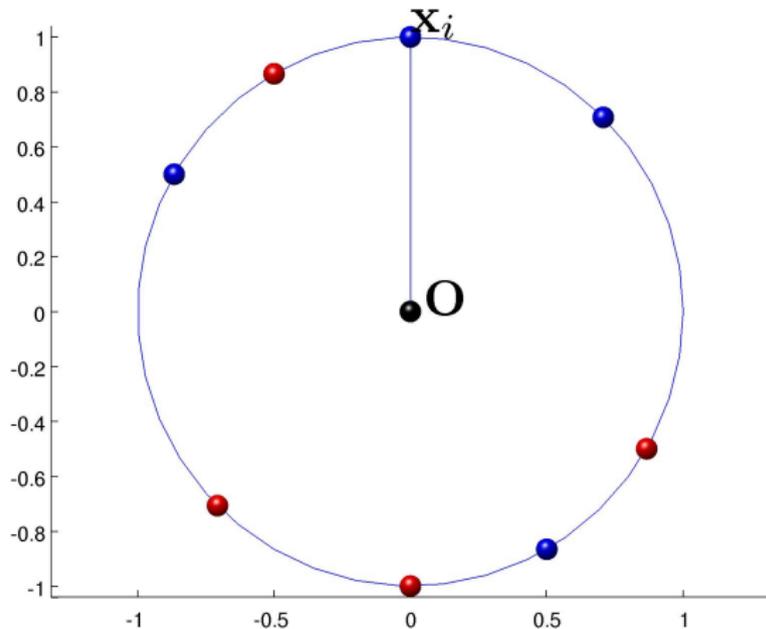
- Outliers

$$\begin{aligned} \min & \| \mathbf{a} \|_1 + \lambda \| \mathbf{e} \|_1 \\ \text{s.t.} & \\ & \mathbf{x}_i = \mathbf{X}_{-i} \mathbf{a} + \mathbf{e} \\ & a_j = 0 \end{aligned}$$

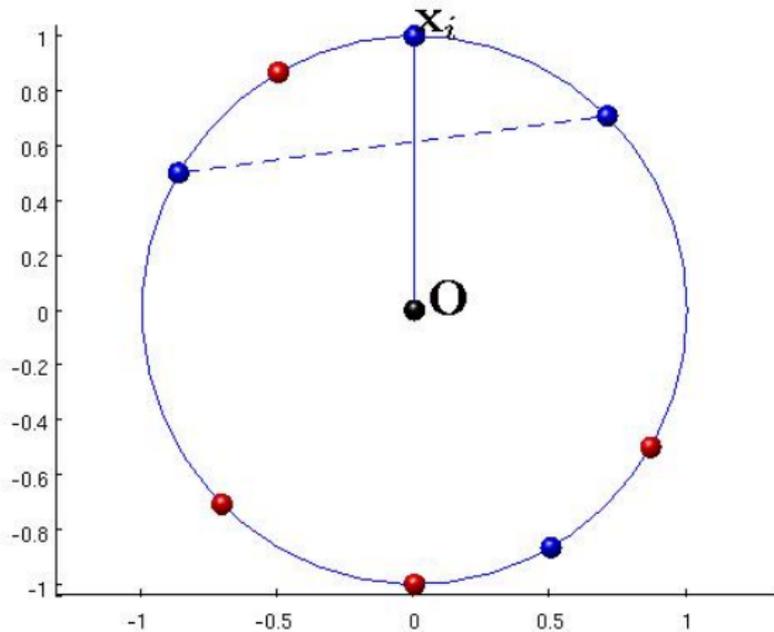
- Noise and Outliers
- Extension to manifolds
- **Graph Connectivity** in each subspace

- 1 Subspace Clustering
- 2 Sparse Subspace Clustering
- 3 Graph Connectivity**
 - Example
 - Preparation
 - 2D Case
 - 3D Case
 - N-D Case
- 4 Conclusion

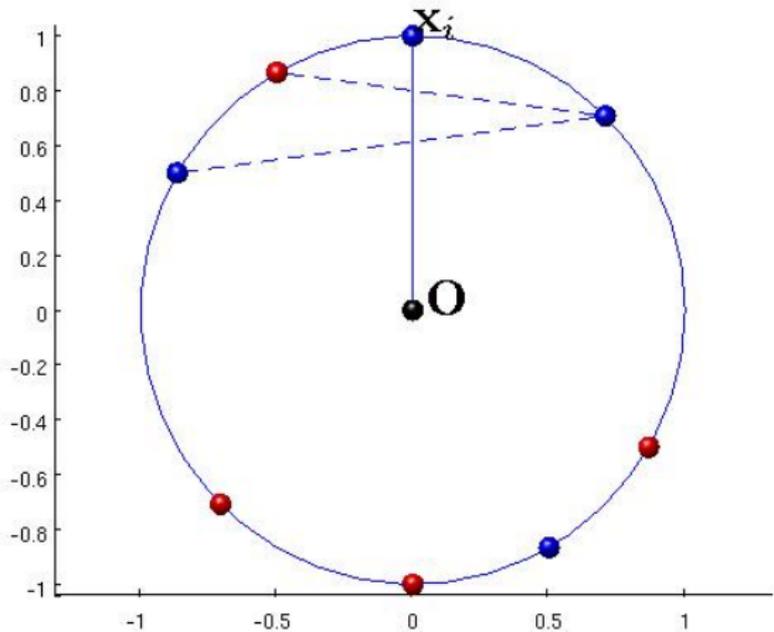
A Simple Example



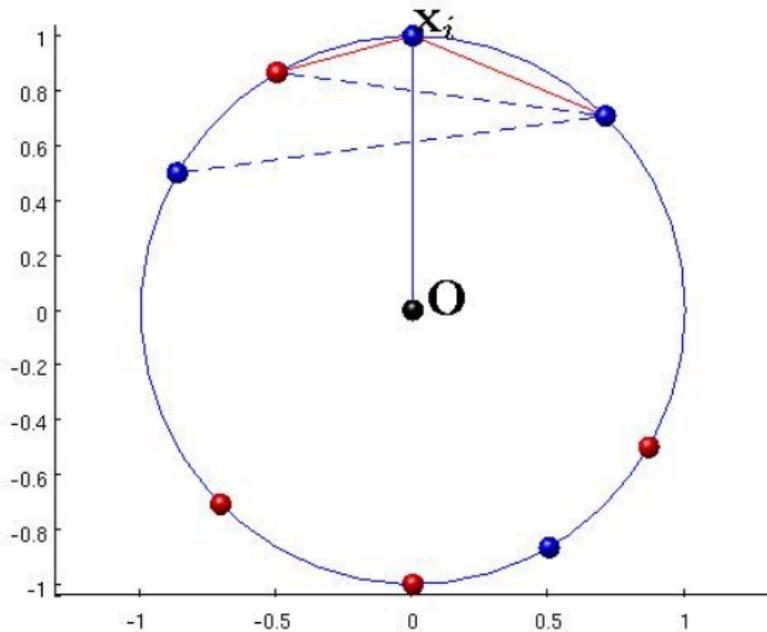
A Simple Example



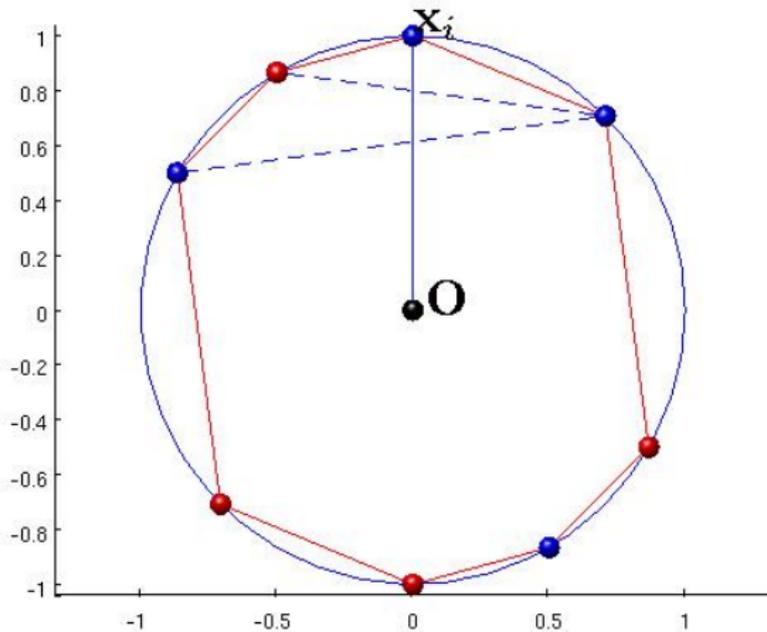
A Simple Example



A Simple Example



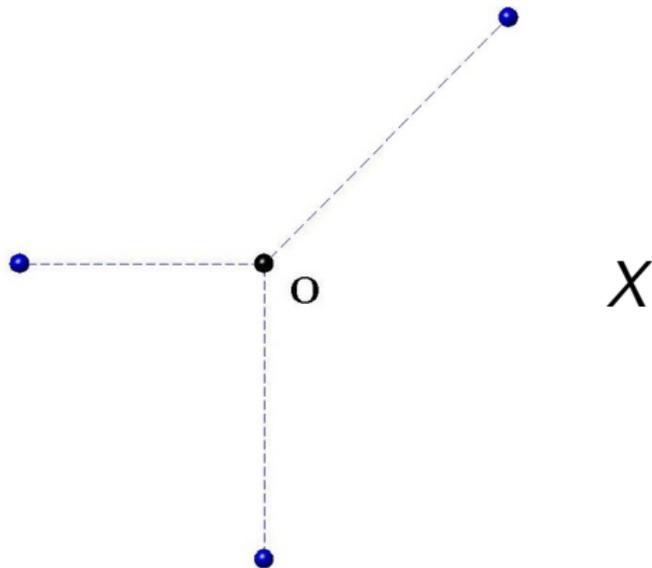
A Simple Example



Preparation

Adding Negative Points

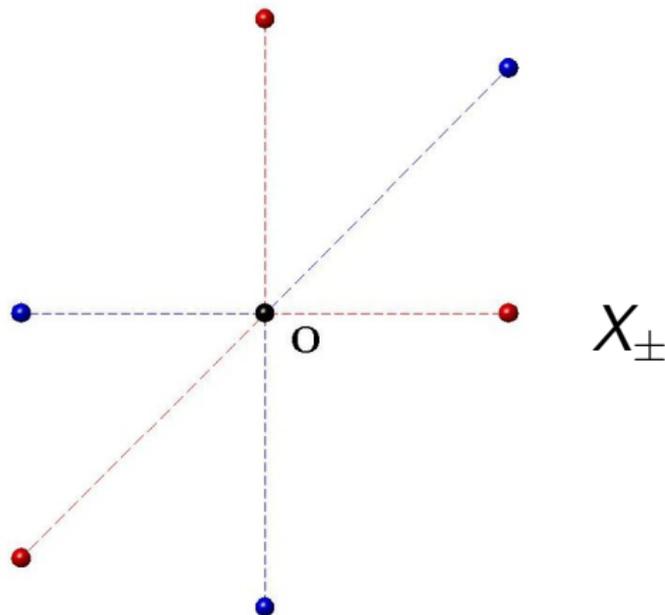
• $\mathbf{x}_i = \sum_{j \neq i} a_j \mathbf{x}_j$



Preparation

Adding Negative Points

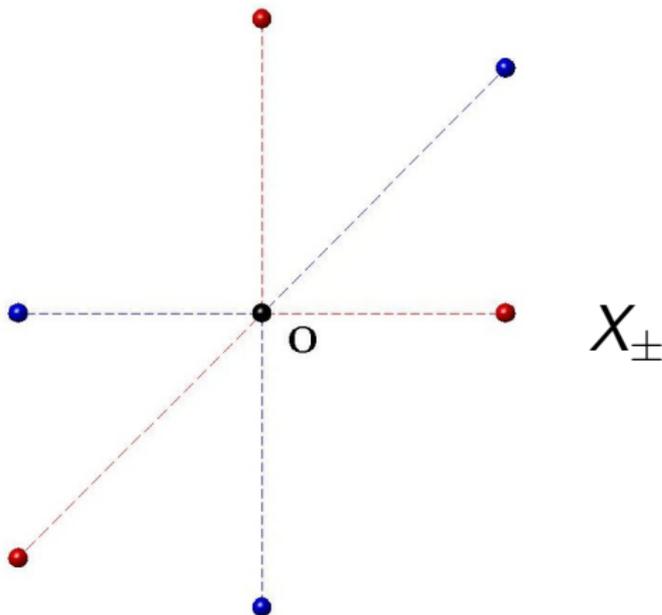
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Preparation

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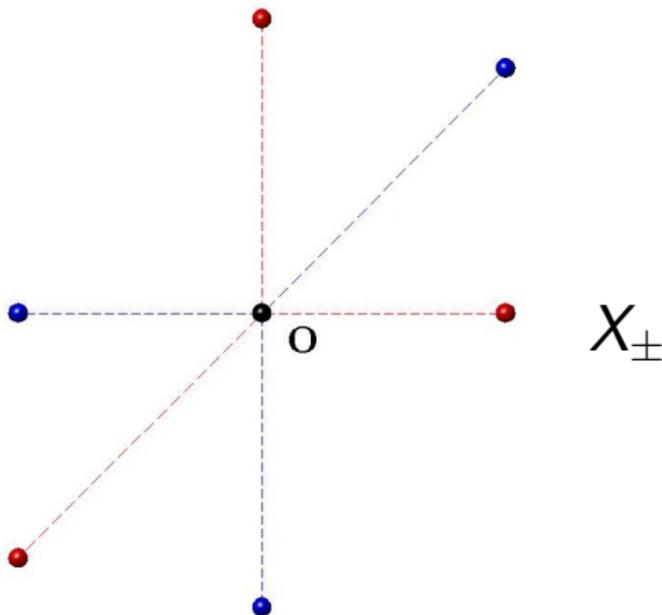
- $\mathbf{x}_i = \sum_{j \neq i} a_j \mathbf{x}_j$
- $a_j \mathbf{x}_j = -a_j(-\mathbf{x}_j)$



Preparation

Adding Negative Points

- $\mathbf{x}_i = \sum_{j \neq i} a_j \mathbf{x}_j$
- $a_j \mathbf{x}_j = -a_j (-\mathbf{x}_j)$
- $a_j \geq 0$



$$\mathbf{a}^i = \arg \min \|\mathbf{a}\|_1$$

s.t.

$$\mathbf{x}_i = \mathbf{X} \mathbf{a},$$

$$a_i = 0$$

$$\mathbf{a}^i = \arg \min \mathbf{1}^T \mathbf{a}$$

s.t.

$$\mathbf{x}_i = \mathbf{X}_{\pm} \mathbf{a},$$

$$\mathbf{a} \succeq \mathbf{0},$$

$$a_i = 0$$

- L^1 minimization and geometry of polytopes *[David Donoho 2005]*

- L^1 minimization and geometry of polytopes [David Donoho 2005]

$$x_i = \mathbf{X}_{-i} \mathbf{b} = \mathbf{X}_{-i} \frac{\mathbf{b}}{\|\mathbf{b}\|_1} \cdot \|\mathbf{b}\|_1$$

- L^1 minimization and geometry of polytopes [David Donoho 2005]

$$\mathbf{x}_i = \mathbf{X}_{-i} \mathbf{b} = \mathbf{X}_{-i} \frac{\mathbf{b}}{\|\mathbf{b}\|_1} \cdot \|\mathbf{b}\|_1$$

$$\mathbf{x}_i = \mathbf{X}_{-i} \mathbf{p} \alpha,$$

where

$$\|\mathbf{p}\| = 1, \mathbf{p} \succeq \mathbf{0}$$

α : to be minimized

- L^1 minimization and geometry of polytopes [David Donoho 2005]

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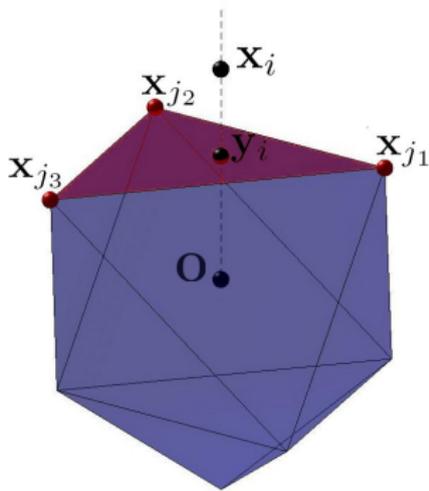
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$$\|\mathbf{p}\| = 1, \mathbf{p} \succeq \mathbf{0}$$

α : to be minimized

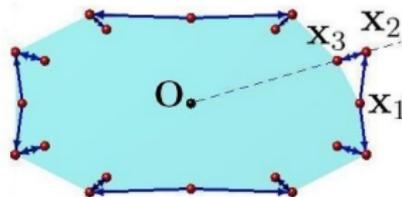
$\mathbf{X}_{-i} \mathbf{p}$: convex hull of \mathbf{X}_{-i}

$$\text{maximize } \beta \quad \text{s.t.} \quad \beta \mathbf{x}_i \in \text{hull}(X_{-i})$$

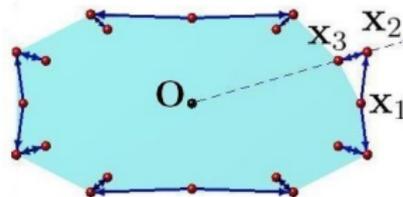


- Indivisible subspaces?

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- Degenerate Cases



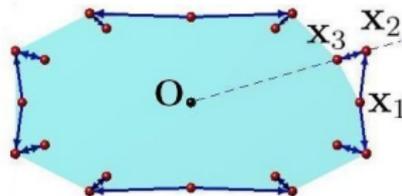
- Indivisible subspaces?
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Assumption

No d points lie in a $(d-1)$ -dimensional subspace.

- Indivisible subspaces?
- Degenerate Cases

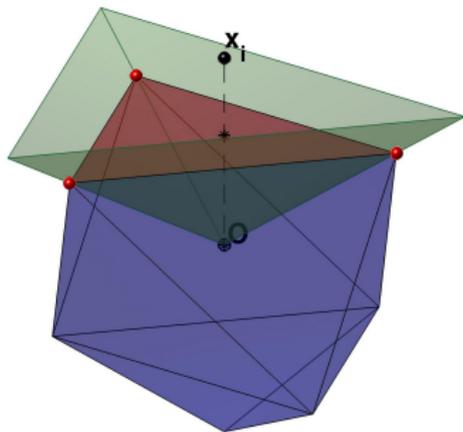


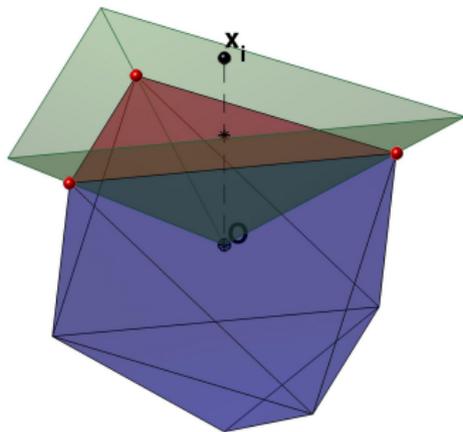
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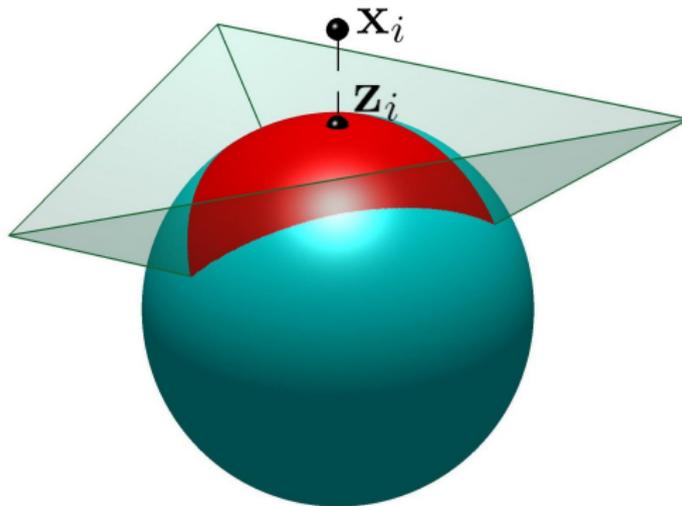
*The facet of each **polytope**(X_{-i}) on which y_i lies has exactly d points on it.*

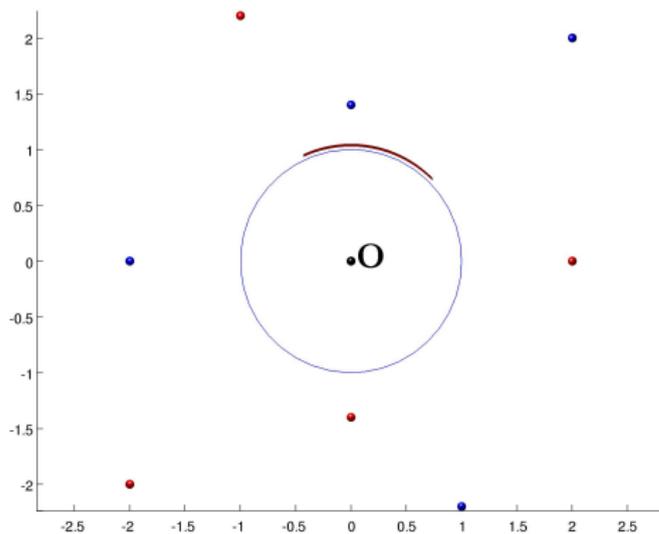


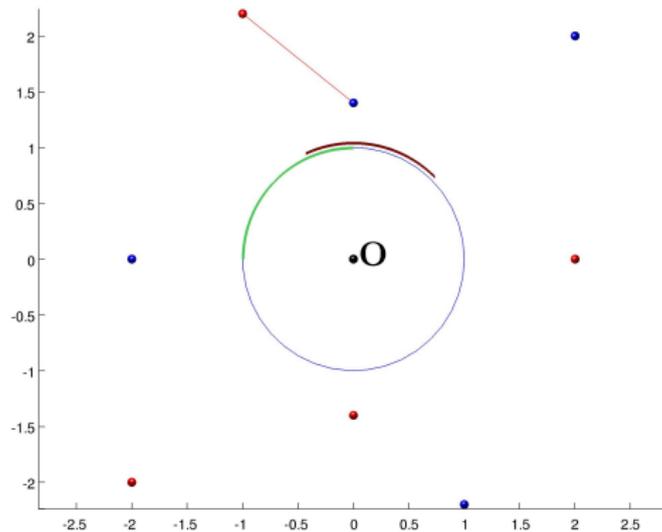


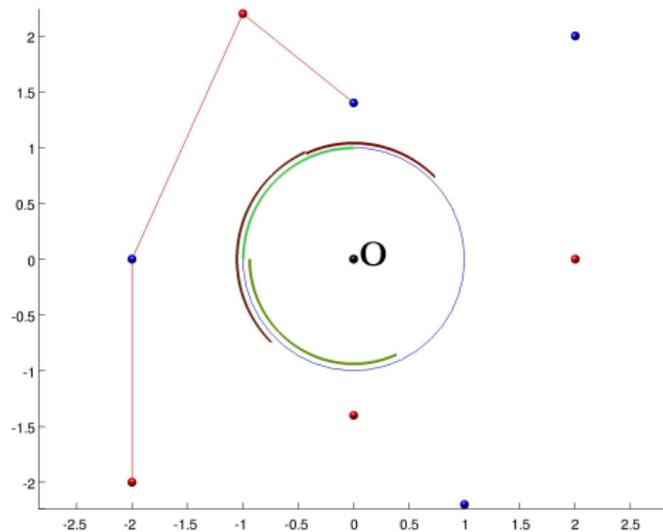
Theorem

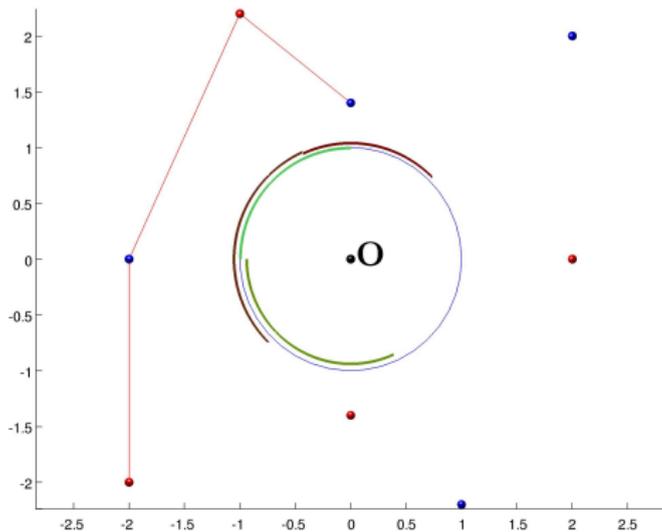
Two points are neighbours iff their neighbourhood cones strictly intersect.

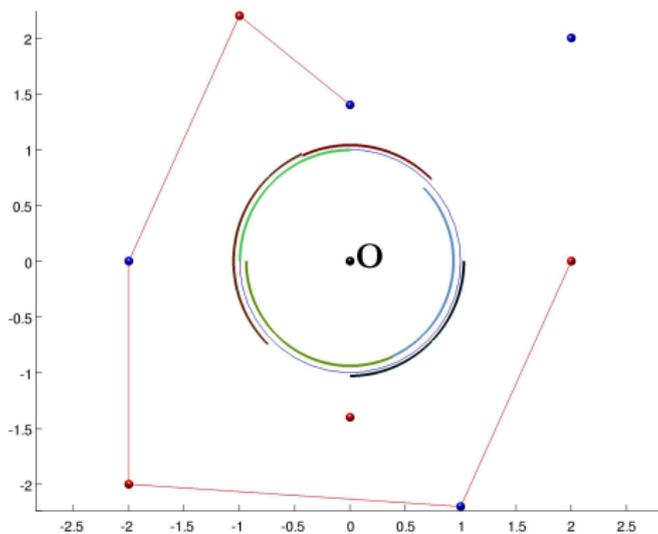


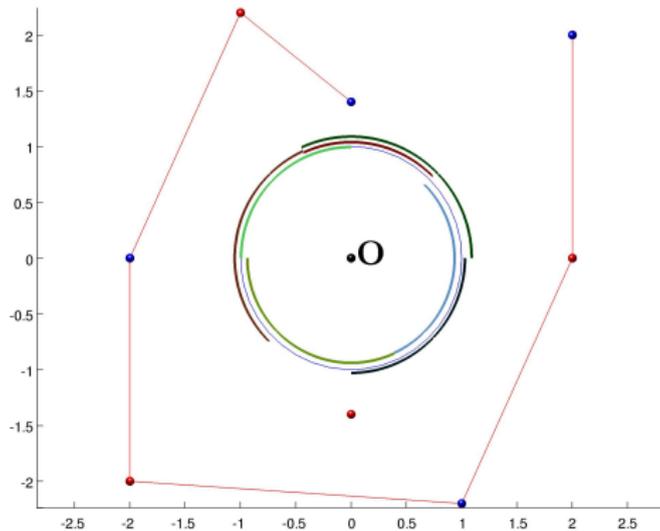


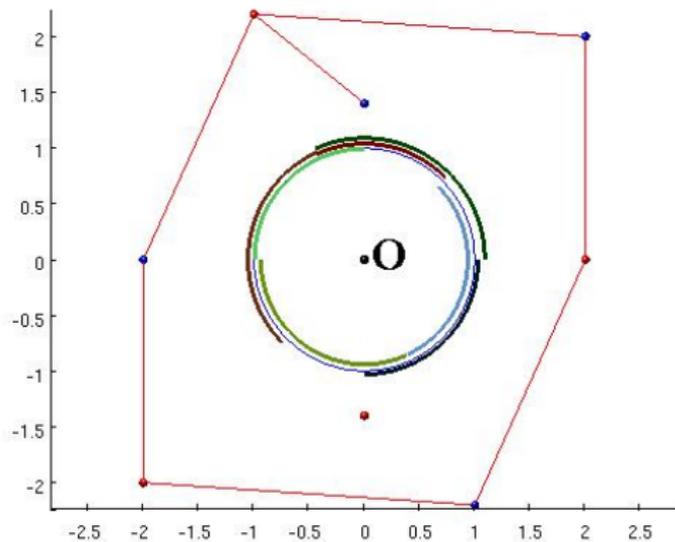


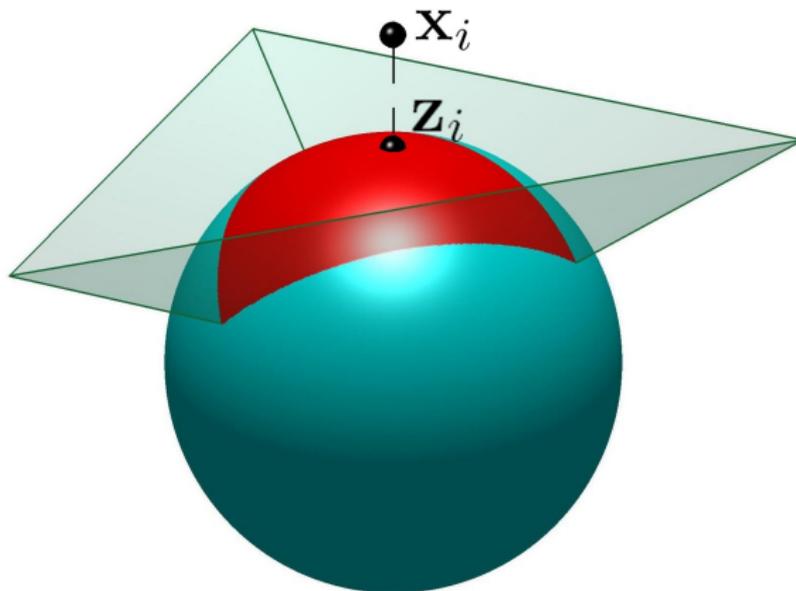


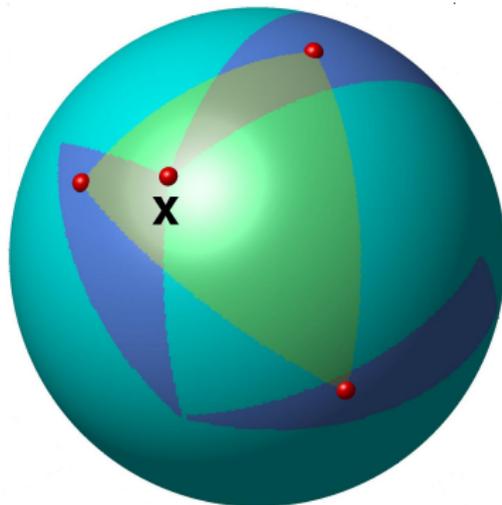


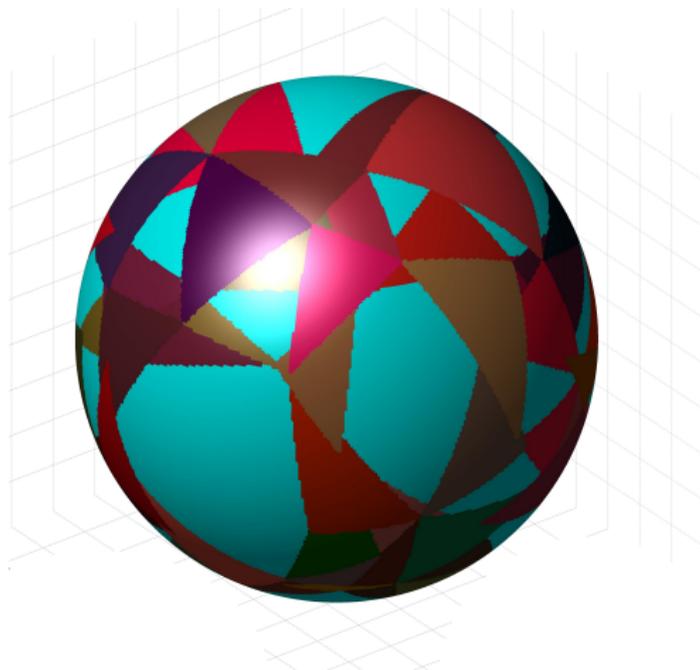


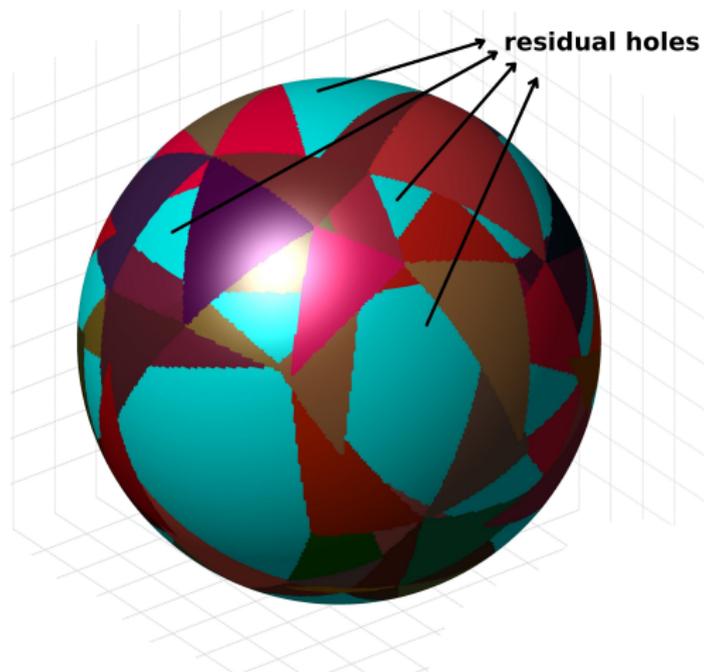


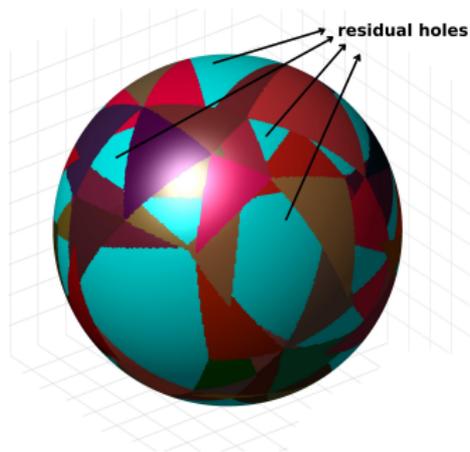












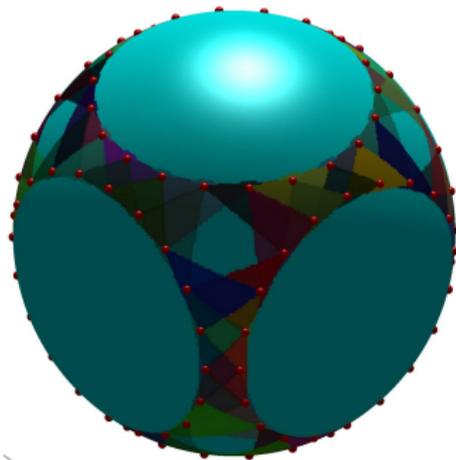
- **residual holes** for one connected component
 - Topologically **open disks**.
 - Area: **< half-sphere**. (\Leftarrow Gauss-Bonnet Theorem)

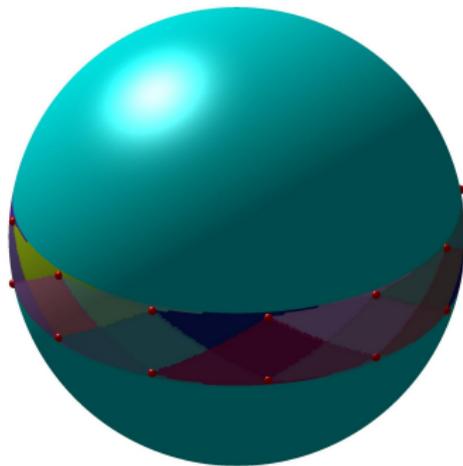
What about $d \geq 4$?



- No residual holes
- Search for counterexamples

Observations from 3D





- Data around two great circles:

- $[\cos \theta, \sin \theta, 0, 0]^T$
- $[0, 0, \cos \theta, \sin \theta]^T$

- $X_{\pm} = X_{C_1} \cup X_{C_2}$

- $X_{C_1} : [\cos \frac{k\pi}{m}, \sin \frac{k\pi}{m}, \pm \delta, \pm \delta]^T$
- $X_{C_2} : [\pm \delta, \pm \delta, \cos \frac{k\pi}{m}, \sin \frac{k\pi}{m}]^T$



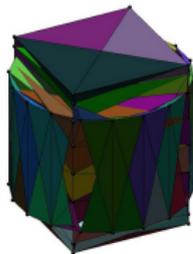
(a) $\delta = \Delta + \epsilon$



(b) $\delta = \Delta + \epsilon$



(c) $\delta = \frac{\sqrt{2}}{2} - \Delta - \epsilon$

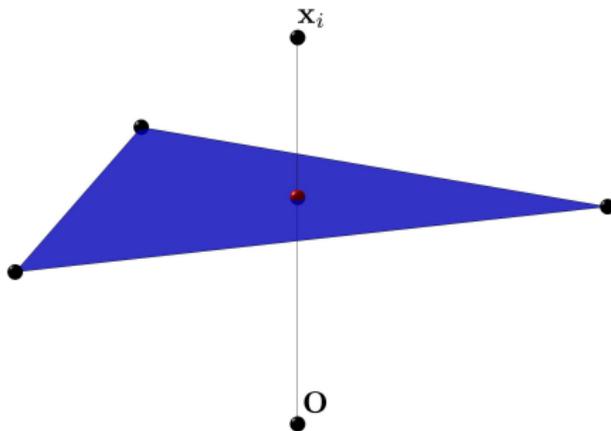


(d) $\delta = \frac{\sqrt{2}}{2} - \Delta + \epsilon$

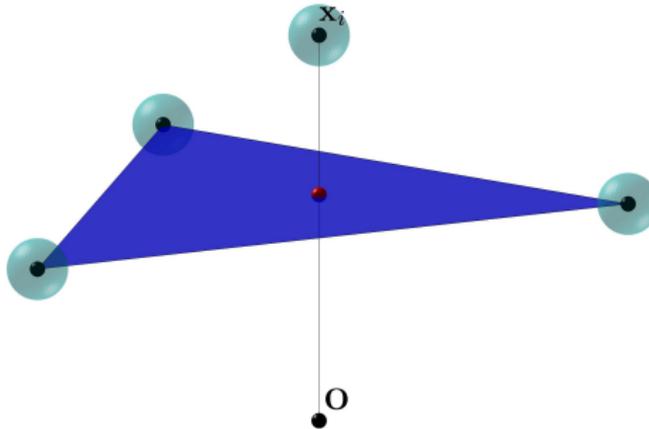
- Disconnected for:
 - $\delta \in (\Delta, \frac{\sqrt{2}}{2} - \Delta)$
 - $\Delta = f(m)$

Figure: Orthographic projection to 3D

- In the counterexample $\text{ray}(x_i)$ hits the interior of the facet.



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- The importance of Subspace Clustering

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- For $d = 2$ and $d = 3$ it will not fail,

- The importance of Subspace Clustering
- Advantages of Sparse Subspace Clustering
- For $d = 2$ and $d = 3$ it will not fail,
- Caution must be taken for $d \geq 4$,
 - A post processing stage, etc.

Thanks

Questions?



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